

Some adventures with compile time evaluations



Mohd Furquan

mohdfurquan1@zhcet.ac.in

mfurquan@iitk.ac.in



**ALIGARH MUSLIM
UNIVERSITY**

IIT KANPUR
Indian Institute of Technology Kanpur



compile time evaluation aka const expr



- Speed-up: calculate once, use ever after!
 - A fraction of computations in a complex program can be evaluated at compile time.
- Forces referential transparency: no state at compile time!
 - Closer to functional paradigm
 - Easier debugging and proving correctness
 - Thread safety
- Language support
 - “they (*languages*) require constant expressions to be written in an impoverished language with minimal support from the type system; this is tedious and error-prone.”
 - Reis & Stroustrup (2010)*

*General constant expressions for system programming languages, 2010 ACM symposium of Applied Computing.

Two approaches: Rosetta code's C++ examples



Templates/Macro style

```
#include <iostream>

template<int i> struct Fac
{
    static const int result = i * Fac<i-1>::result;
};

template<> struct Fac<1>
{
    static const int result = 1;
};

int main()
{
    std::cout << "10! = " << Fac<10>::result << "\n";
    return 0;
}
```

Function composition style

```
#include <stdio.h>

constexpr int factorial(int n) {
    return n ? (n * factorial(n - 1)) : 1;
}

constexpr int f10 = factorial(10);

int main() {
    printf("%d\n", f10);
    return 0;
}
```

Poor man's approach: Rosetta code's Fortran example



```

program test

  implicit none
  integer,parameter :: t = 10*9*8*7*6*5*4*3*2 !computed at compile time

  write(*,*) t !write the value the console.

end program test

```

**Legit., but
useless!**
**Is it so bad
for Fortran?**

Founder's* vision:

Def Innerproduct

$\equiv (\text{Insert } +) \circ (\text{ApplyToAll } \times) \circ \text{Transpose}$

Or, in abbreviated form:

Def IP $\equiv (/+) \circ (\alpha \times) \circ \text{Trans.}$

IP: $\langle\langle 1,2,3 \rangle, \langle 6,5,4 \rangle\rangle =$

Definition of IP

$\Rightarrow (/+) \circ (\alpha \times) \circ \text{Trans: } \langle\langle 1,2,3 \rangle, \langle 6,5,4 \rangle\rangle$

Effect of composition, \circ

$\Rightarrow (/+):((\alpha \times):(\text{Trans: } \langle\langle 1,2,3 \rangle, \langle 6,5,4 \rangle\rangle))$

Applying Transpose

$\Rightarrow (/+):((\alpha \times): \langle\langle 1,6 \rangle, \langle 2,5 \rangle, \langle 3,4 \rangle\rangle)$

Effect of ApplyToAll, α

$\Rightarrow (/+): \langle \times: \langle 1,6 \rangle, \times: \langle 2,5 \rangle, \times: \langle 3,4 \rangle \rangle$

Applying \times

$\Rightarrow (/+): \langle 6, 10, 12 \rangle$

Effect of Insert, /

$\Rightarrow +: \langle 6, +: \langle 10, 12 \rangle \rangle$

Applying +

$\Rightarrow +: \langle 6, 22 \rangle$

Applying + again

$\Rightarrow 28$

*Backus, John, Can programming be liberated from the von Neumann style?: a functional style and its algebra of programs. ACM Turing Award Lecture, 1977.

Translating into Modern Fortran:



```
program InnerProduct
  implicit none
  integer,parameter :: input=[[1,2,3],[6,5,4]]

  write(*,*) SUM(PRODUCT(input,2))  !should be computed at compile time
end program InnerProduct
```

Approach:

- Default constructor
- Composition of intrinsic functions

Improved Rosetta code's example:

```
program factorial
  implicit none

  integer :: i
  integer,parameter :: N, list = [(i, i = 1,N)]

  write(*,*) PRODUCT(list)  !should be computed at compile time
end program factorial
```

Not as general
as

**C++ constexpr
functions,
useful still!**

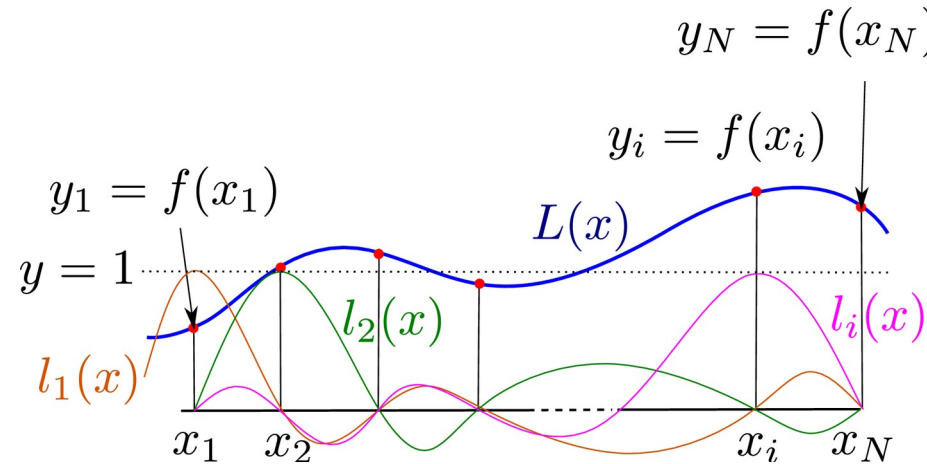
A complex example:

Lagrange polynomials & derivatives



Lagrange basis:

$$l_i(x) = \prod_{\substack{k=1 \\ k \neq i}}^{k=N} \left(\frac{x - x_k}{x_i - x_k} \right)$$
$$l'_i(x) = \frac{\sum_{j=1}^N \left(\prod_{\substack{k=1 \\ k \neq i, j}}^N (x - x_k) \right)}{\prod_{\substack{k=1 \\ k \neq i}}^N (x_i - x_k)} = \left(\sum_{\substack{k=1 \\ k \neq i}}^N \frac{1}{x - x_k} \right) l_i(x)$$



Lagrange polynomials:

$$L(x) = \sum_{i=1}^N y_i l_i(x)$$

Task: Evaluate l_i and l_i' at selected M values of x, using N points (x_i 's)

Selecting x from table:

Choosing quadrature rule



Table of 1D quadrature points:

```
! ADJUSTABLE PARAMETER
integer      ,parameter      :: M

! 1D Gaussian quadrature
integer      ,parameter,private :: maxint = 6
real(kind=rp),parameter,private ::           &
xg(maxint) = [                               & ! coordinates
  0._rp,                                       & ! 1 point
 -1._rp/sqrt(3._rp),1._rp/sqrt(3._rp),       & ! 2 point
 -sqrt(3._rp/5._rp),0._rp,sqrt(3._rp/5._rp)] ! 3 point

! selecting quadrature points
integer      ,parameter      ::           &
qi   = (M - 1)*M/2 + 1,         &
qf   = (M + 1)*M/2
```

Approach:

- Use parameters as 'keys' and 'levers'.

- Declare private parameters for intermediate calculations.

Required N values of x: **xg(qi:qf)**

Computing $l_i(x)$ via x_i & x_k : SPREAD all over!



```
! ADJUSTABLE PARAMETER
```

```
integer,parameter :: N
```

```
integer,parameter,private :: i, j, k, p, q, r
```

```
logical,parameter,private,dimension(M,N,N) :: mq = spread(reshape([(i/=j, i = 1,N), j = 1,N]), [N,N]),1,M) &
```

```
real(kind=rp),parameter,private,dimension(N) :: xn = [(-1+2*i/(N-1), i=0,N-1)] &
```

```
real(kind=rp),parameter,private,dimension(nqd,nen,nen) :: x = spread(spread(xg(qi:qf),2,N),3,N), & & xk = spread(spread(xn,1,N),1,M), & xi = spread(spread(xn,1,M),3,N)
```

```
real(kind=rp),parameter,private,dimension(M,N) :: li = product(x - xk,3,mq)/product(xi - xk,3,mq) &
```

Approach:

- Replace loops by vector operations.
- Use SPREAD to generate unrolled loops.

$$l_i(x) = \prod_{\substack{k=1 \\ k \neq i}}^{k=N} \left(\frac{x - x_k}{x_i - x_k} \right)$$

Computing $l_i'(x)$:

TWO approaches!



Approach 1:

- Simple, no new definition is required.
- **Zero denominator will cause problem!**

```
real(kind=rp),parameter          ,dimension(M,N) ::  
  li_x = sum(1./(x - xk),3,mq)*li
```

$$l_i'(x) = \left(\sum_{\substack{k=1 \\ k \neq i}}^N \frac{1}{x - x_k} \right) l_i(x)$$

Approach 2:

- Requires definition of an additional mask.
- Complex computation.

```
logical          ,parameter,private,dimension(M,N,N,N) ::  
  cq = spread(reshape([(((i/=j.AND.(i-j/=k.AND.i-j/=k-N),  
    i = 1,N), j = 1,N), k = 1,N-1)],  
    [N,N,N-1]),1,M)  
  
real(kind=rp),parameter          ,dimension(M,N) ::  
  li_x = sum(product(spread(x - xk,4,N-1),3,cq),3)  
    / product(          xi - xk,3,mq)
```

$$l_i'(x) = \frac{\sum_{j=1}^N \left(\prod_{\substack{k=1 \\ k \neq i,j}}^N (x - x_k) \right)}{\prod_{\substack{k=1 \\ k \neq i}}^N (x_i - x_k)}$$

Multidimensional polynomials: via tensor products!



$$m_{ijk}(x, y, z) = l_i(x)l_j(y)l_k(z)$$

```
real(kind=rp),parameter,dimension(M**3,N**3) ::      &
  mijk = reshape([(((li(i,p)*li(j,q)*li(k,r),        &
                    i = 1,M), j = 1,M), k = 1,M),    &
                    p = 1,N), q = 1,N), r = 1,N)],   &
                [M**3,N**3])
```

$$m_{ijk,x}(x, y, z) = l'_i(x)l_j(y)l_k(z)$$

```
real(kind=rp),parameter,dimension(3,N**3,M**3) ::  &
  mijk_x = reshape([(((li_x(i,p)*li (j,q)*li (k,r),  &
                      li (i,p)*li_x(j,q)*li (k,r),  &
                      li (i,p)*li (j,q)*li_x(k,r),   &
                    p = 1,N), q = 1,N), r = 1,N),    &
                    i = 1,M), j = 1,M), k = 1,M)],   &
                [3,N**3,M**3])
```

Approach:

- Implied loops
and RESHAPE to
accomplish tensor
products.

Thank you!

Questions?

